# When Is Linear Recursion Non-negative 

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What about $u_{3}=(2 \sqrt{5}+3) u_{2}-3(2 \sqrt{5}+3) u_{1}-27 u_{0}$ ?

## Definition

Let $\mathbf{u}=\left\{u_{n}\right\}_{n \geq 0}$ be a linear recurrence sequence (LRS) defined by

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## Definition

We say $\mathbf{u}$ is:
( positive if $u_{n} \geq 0$ for all $n \geq 0$
( ultimately positive if $\exists N$ so $u_{n} \geq 0$ for all $n \geq N$

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- this is called the Skolem's problem
- question 2 is solved
- question 3 and 4 are still open for large depth


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- In biology we have what's called L-system, which was originated from simulating the development of multicellular organisms
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O In computer science we have verification of lienar automata We want to know whether those sequences will be ultimately positive or not because some models would not have real-life meaning if the values are negative

## Definitions

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- The roots of $f_{\mathbf{u}}(x)$ are called characteristic roots of $\mathbf{u}$
- The dominant roots of $\mathbf{u}$ are the roots with maximum modulus.


## Theorem

## Theorem 2

Let u be a LRS, then

$$
u_{n}=p_{1}(n) \gamma_{1}^{n}+\ldots+p_{k}(n) \gamma_{k}^{n}
$$

where $p_{i}(x) \in \mathbb{C}[x]$ are polynomials, and $\gamma_{i}$ are the characteristic roots.

## Proof: Simple Case

We first deal with $f_{\mathbf{u}}(x)=\left(x-\gamma_{1}\right) \ldots\left(x-\gamma_{k}\right)$, i.e. $f_{\mathbf{u}}$ has $k$ distinct complex roots.

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Define

$$
C_{\mathbf{u}}:=\left[\begin{array}{ccccccc}
a_{1} & a_{2} & a_{3} & a_{4} & \ldots & a_{k-1} & a_{k} \\
1 & 0 & 0 & 0 & \ldots & 0 & 0 \\
0 & 1 & 0 & 0 & \ldots & 0 & 0 \\
\vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\
0 & 0 & 0 & 0 & \ldots & 1 & 0
\end{array}\right]
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Then we see

$$
\left[\begin{array}{c}
u_{n+k} \\
u_{n+k-1} \\
\vdots \\
u_{n+1}
\end{array}\right]=C_{\mathbf{u}}\left[\begin{array}{c}
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$$
\left[\begin{array}{c}
u_{n} \\
u_{n-1} \\
\vdots \\
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\end{array}\right]=P D^{n} P^{-1}\left[\begin{array}{c}
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## Remark

Since $P, D$ are all matrices, we see in this case $p_{i}$ in our theorem are all constant polynomials.

## Example

Take our favorite sequence $F_{n}=F_{n-1}+F_{n-2}$ with $F_{0}=F_{1}=1$. Then

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C_{F}=\left[\begin{array}{ll}
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The diagonalization is given by

$$
C_{F}=\left[\begin{array}{cc}
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2 & 2
\end{array}\right]\left[\begin{array}{cc}
\phi_{1} & 0 \\
0 & \phi_{2}
\end{array}\right]\left[\begin{array}{cc}
\frac{1}{2 \sqrt{5}} & -\frac{1-\sqrt{5}}{4 \sqrt{5}} \\
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where $\phi_{i}=\frac{1+(-1)^{i} \sqrt{5}}{2}$

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$$
F_{n}=\frac{1}{\sqrt{5}} \phi_{2}^{n}-\frac{1}{\sqrt{5}} \phi_{1}^{n}
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## Proof: General Case

In this case diagonalization is no longer helpful, as in this case one can actually show $C_{\mathbf{u}}$ is diagonalizable iff $f_{\mathbf{u}}$ has $k$ distinct roots.

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- Assume $n \geq k$, then we see

$$
\left\{\begin{array}{l}
U(x)=\ldots+u_{n} x^{n} \\
x U(x)=\ldots+u_{n-1} x^{n} \\
x^{2} U(x)=\ldots+u_{n-2} x^{n} \\
\vdots
\end{array}\right.
$$

## Proof: General Case

- Thus we see

$$
\left(1-a_{1} x-a_{1} x^{2}-\ldots-a_{k} x^{k}\right) U(x)=\ldots+\left(u_{n}-a_{1} u_{n-1}-\ldots\right) x^{n}+\ldots
$$

- In other word, the RHS cannot contain terms with degree higher than $k-1$. Denote this by $G(x)$.
O Thus

$$
\left(x^{k} f_{\mathbf{u}}(1 / x)\right) U(x)=G(x) \Rightarrow U(x)=\frac{G(x)}{x^{k} f_{\mathbf{u}}(1 / x)}
$$

- Therefore, $u_{n}$ is exactly the $n$th coefficient of the Taylor expansion of $G(x) /\left(x^{k} f_{\mathbf{u}}(1 / x)\right)$ around $x=0$, but one can verify that this gives the desired closed form


## Examples

## Example

- If $u_{n}=2 u_{n-1}-u_{n-2}$ and $u_{0}=3, u_{1}=1$, then $f_{u}(x)=(x-1)^{2}$ and

$$
u_{n}=p_{1}(n) \cdot 1+p_{2}(n) \cdot 1, \quad p_{1}(n)=3, p_{2}(n)=-2 n
$$

- If $u_{n}=2 u_{n-1}+u_{n-2}$ with $\left(u_{0}, u_{1}\right)=(3,1)$ then

$$
u_{n}=\frac{1}{2}\left((3+\sqrt{2}) \gamma_{1}^{n}-(\sqrt{2}-3) \gamma_{2}^{n}\right)
$$

with $\gamma_{1}=1-\sqrt{2}$ and $\gamma_{2}=1+\sqrt{2}$.

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Ultimate positivity problem for simple LRS is decidable

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Next, we will try to give a basic idea of how to prove a statement like this

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## Definition 5

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The number $\phi=\frac{1+\sqrt{5}}{2}$ is algebraic as $f(x)=x^{2}-x-1$ vanishes $\phi$. On the other hand $\pi$ is not algebraic (this is non-trivial).

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Let $\alpha$ be algebraic, then:

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O the collection of algebraic integers is denoted by $\mathcal{O}$

## Number Fields

## Definition 8

A number field $K$ is a field $K$ such that $\mathbb{Q} \subseteq K$ and $K$ is finite dimensional vector space over $\mathbb{Q}$. The ring of integers for $K$ is defined by $\mathcal{O}_{K}:=\mathcal{O} \cap K$.

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## Example 9

Consider the $\mathbb{Q}$-vector space $R=\mathbb{Q}[x]$ and quotient space $K:=R / I$ where $I$ is the subspace spanned by vector $\left\{\left(x^{2}-5\right) \cdot f: f \in R\right\}$. In other word,

$$
I=\left\{f \in R:\left(x^{2}-5\right) \mid f\right\}
$$

Then $K$ is a number field, and it can be written as $\mathbb{Q}(\sqrt{5})$, i.e. $K$ is isomorphic to the vector space spanned by two elements, 1 and $\sqrt{5}$ by the isomorphism $1 \mapsto 1$ and $x \mapsto \sqrt{5}$.

## Ideals

## Definition 10

A subset $I \subseteq \mathcal{O}_{K}$ is an ideal if $a \in \mathcal{O}_{K}$ and $b \in I$ then $a b \in I$, and $a, b \in I$ implies $a+b \in I$

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Let $P$ be an ideal of $\mathcal{O}_{K}$, then $P$ is prime if $a b \in P$ then $a \in P$ or $b \in P$

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## Example 12

Consider the number field $K=\mathbb{Q}(\sqrt{-5})=\mathbb{Q}[x] /\left(x^{2}+5\right)$. One can show $\mathcal{O}_{K}=\mathbb{Z}[\sqrt{-5}]$ in this case, and the ideal 2 is not prime. To see this, note $2=(1+\sqrt{-5})(1-\sqrt{-5})$ while none of those two elements lies in the ideal

$$
\text { (2) }:=\{2 a+2 b \sqrt{-5}: a, b \in \mathbb{Z}\}
$$

## Ideals

## Example 13

Let $K=\mathbb{Q}(\sqrt{-5})$ and consider the ideal (2). Then we get

$$
(2)=(1+\sqrt{-5}, 2)^{2}
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where $(1+\sqrt{-5}, 2)$ is a prime ideal.

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Every ideal $I \subseteq K$ admits a unique prime factorization, i.e. there exists prime ideals $P_{i}$ so $I=\prod P_{i}^{k_{i}}$.

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## Remark

Note here product of ideals $I J$ is defined by

$$
I J:=\{a b: a \in I, b \in J\}
$$

## S-Units

## Definition 15

Let $K$ be a number field, and $S$ a finite set of prime ideals in $\mathcal{O}_{K}$. Then $\alpha \in \mathcal{O}_{K}$ is a $S$-unit if the principal ideal $(\alpha)$ is a product of prime ideals in $S$.

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## S-Units

## Definition 15

Let $K$ be a number field, and $S$ a finite set of prime ideals in $\mathcal{O}_{K}$. Then $\alpha \in \mathcal{O}_{K}$ is a $S$-unit if the principal ideal $(\alpha)$ is a product of prime ideals in $S$.

## Theorem 16

Let $K$ be a number field, $m$ positive integer, $S$ a finite set of primes in $\mathcal{O}_{K}$. Then for every $\epsilon>0$ there exists constant $C$, depending on $m, K, S$ and $\epsilon$, such that, for any $S$-units $x_{1}, \ldots, x_{m}$ with no proper subsum equal 0 , we have

$$
\left|x_{1}+\ldots+x_{m}\right| \geq C X Y^{-\epsilon}
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where $X, Y$ are computable constants only depending on $K$ and $x_{i}$.

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## Remark

The proof of the above theorem uses a very vast generalization of what's called Roth's theorem, which is a field medal result. For a pointer, Roth's theorem admits a vast generalization called Schmidt's subspace theorem, and this theorem uses the p-adic version of subspace theorem.

## Result

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Let $\mathbf{u}$ be non-degenerate and simple. Then, there exists a function $F: \mathbb{T}^{s} \rightarrow \mathbb{R}$, depending on $\mathbf{u}$, so that $\mathbf{u}$ is ultimately positive iff $F(\mathbf{z}) \geq 0$ for all $\mathbf{z} \in T(\mathbf{a}) \subseteq \mathbb{T}^{s}$.

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Here $\mathbb{T}=\{z \in \mathbb{C}:|z|=1\}$ and $\mathbb{T}^{s}$ is $s$ copy of $\mathbb{T}$

## Definitions

Suppose $\mathbf{u}$ is simple and non-degenerate, with dominant roots

$$
\rho, \gamma_{1}, \ldots, \gamma_{s}, \overline{\gamma_{1}}, \ldots, \overline{\gamma_{s}}
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where $\rho$ is real and positive.

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Then,

$$
u_{n}=b \rho^{n}+\sum_{i=1}^{s} c_{i} \gamma_{i}^{n}+{\overline{c_{i} \gamma_{i}}}^{n}+r(n)
$$

where $r(n)$ is relatively small.

## Definitions

Now set $\lambda_{i}=\gamma_{i} / \rho$, then

$$
u_{n}=\rho^{n} F\left(\lambda_{1}^{n}, \ldots, \lambda_{s}^{n}\right)+r(n)
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where

$$
F\left(z_{1}, \ldots, z_{s}\right)=b+c_{1} z_{1}+\ldots+c_{s} z_{s}+\overline{c_{1} z_{1}}+\ldots+\overline{c_{s} z_{s}}
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## Definition

Let $\mathbf{a}=\left(\lambda_{1}, \ldots, \lambda_{s}\right)$ as above, we define

$$
L(\mathbf{a}):=\left\{\left(v_{1}, \ldots, v_{s}\right) \in \mathbb{Z}^{s}: a_{1}^{v_{1}} \ldots a_{s}^{v_{s}}=1\right\} \subseteq\left(\mathbb{Z}^{s},+\right)
$$

$$
T(\mathbf{a}):=\left\{\left(\mu_{1}, \ldots, \mu_{s}\right) \in \mathbb{T}^{s}: \mu_{1}^{v_{1}} \ldots \mu_{s}^{v_{s}}=1 \text { for all } \mathbf{v} \in L(\mathbf{a})\right\}
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\end{gathered}
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This completes our main criterion for ultimate positivity when $\mathbf{u}$ is simple and non-degenerate

Sketch Proof

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## Theorem 18

Let u be non-degenerate and simple. Then, there exists a function $F: \mathbb{T}^{s} \rightarrow \mathbb{R}$, depending on $\mathbf{u}$, so that $\mathbf{u}$ is ultimately positive iff $F(\mathbf{z}) \geq 0$ for all $\mathbf{z} \in T(\mathbf{a}) \subseteq \mathbb{T}^{s}$.

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Now we proceed on give a sketch proof

## Sketch Proof

Let $\rho, \gamma_{1}, \ldots, \gamma_{s}, \overline{\gamma_{1}}, \ldots, \overline{\gamma_{s}}$ be the dominant roots of $f_{\mathbf{u}}$. Then we get

$$
u_{n}=b \rho^{n}+\sum c_{i} \gamma_{i}^{n}+\overline{c_{i} \gamma_{i}^{n}}+r(n)
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where $r(n)=o\left(\rho^{n(1-\epsilon)}\right)$ for some $\epsilon>0$

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Let $K / \mathbb{Q}$ be the number field generated by all the characteristic roots of $\mathbf{u}$, and $S$ the set of prime ideal divisors of the dominant characteristic roots.

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By construction

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## Theorem 19

Let $K$ be a number field, $m$ positive integer, $S$ a finite set of primes in $\mathcal{O}_{K}$. Then for every $\epsilon>0$ there exists constant $C$, depending on $m, K, S$ and $\epsilon$, such that, for any $S$-units $x_{1}, \ldots, x_{m}$ with no proper subsum equal 0 , we have

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where $X, Y$ are computable constants only depending on $K$ and $x_{i}$.
Here $X=C_{1} \rho^{n}, Y=C_{2} \rho^{n}$ for some constant $C_{1}, C_{2}>0$. Thus for all $\epsilon>0$ there exists constant $C$ so

$$
\left|b \rho^{n}+\sum c_{i} \gamma_{i}^{n}+\overline{c_{i} \gamma_{i}^{n}}\right| \geq C \rho^{n(1-\epsilon)}
$$

for all but finitely many $n$

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Now just pick the $\epsilon$ so $r(n)=o\left(\rho^{n(1-\epsilon)}\right)$, we conclude the sum of dominant roots indeed dominant the whole sum for all but finitely many $n$.

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Thus $u_{n}$ ultimately positive iff

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ultimately positive.
But then $b \rho^{n}+\sum c_{i} \gamma_{i}^{n}+\overline{c_{i} \gamma_{i}^{n}}$ can be re-write as $\rho^{n} F\left(\lambda_{1}^{n}, \ldots, \lambda_{s}^{n}\right)$, where the positivity of $F$ can be checked on the torus, which is a computable task.

## General Case

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Hence we get general algorithm as follows:

- Find subsequences $\mathbf{u}_{i}$ for $1 \leq i \leq M$ such that the union of $\mathbf{u}_{i}$ is equal $\mathbf{u}$, where all $\mathbf{u}_{i}$ are non-degenerate
- Verify those $\mathbf{u}_{i}$ individually using the criterion above


## Example 1

Suppose

$$
u_{n}=\frac{3}{2} u_{n-1}+\frac{3}{2} u_{n-2}-u_{n-3}
$$

Then

$$
f_{u}=(x-2)(x+1)\left(x-\frac{1}{2}\right)
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Then

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$$

Hence

$$
\begin{aligned}
u_{n}= & 2^{n}\left(\frac{8}{9} u_{2}+\frac{4}{9} u_{1}-\frac{4}{9} u_{0}\right) \\
& +(-1)^{n}\left(\frac{2}{9} u_{2}-\frac{5}{9} u_{1}+\frac{2}{9} u_{0}\right) \\
& +\frac{-\frac{1}{9} u_{2}+\frac{1}{9} u_{1}+\frac{2}{9} u_{0}}{2^{n}}
\end{aligned}
$$

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## Conclusion

$u_{n}=\frac{3}{2} u_{n-1}+\frac{3}{2} u_{n-2}-u_{n-3}$ is ultimately positive if and only if $\left(u_{2}, u_{1}, u_{0}\right)$ lies in a half-space defined by $(8 / 9,4 / 9,-4 / 9)$.

## Example 2

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u_{n}=(2 \sqrt{5}+3) u_{n-1}-3(2 \sqrt{5}+3) u_{n-2}-27 u_{n-3}
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Then

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\begin{aligned}
f_{\mathbf{u}} & =x^{3}-(2 \sqrt{5}+3) x^{2}+3(2 \sqrt{5}+3) x-27 \\
& =(x-3)\left(x-\gamma_{1}\right)\left(x-\gamma_{2}\right)
\end{aligned}
$$

where $\gamma_{1}, \gamma_{2}$ are conjugate to each other.

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- the numerical instability is mostly coming from the error due to exponents
- to compute $F(\mathbf{z})$ we do not need the exponents!


## Example 2

After some computation, we obtain

$$
\begin{aligned}
F(z) \approx & 18 u_{0}-9 u_{1}+2 u_{2} \\
& +\left(-(9-3 i) u_{0}+(4-4 i) u_{1}-(0.5-1 i) u_{2}\right) z \\
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$$

So

$$
T\left(\lambda_{1}\right)=\left\{z \in \mathbb{C}:|z|=1, z^{0}=1\right\}=\mathbb{T}
$$

## Example 2

Therefore, to see if $\mathbf{u}$ is ultimately positive or not, we just need to find the values of $u_{0}, u_{1}, u_{2}$ such that

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\min (F(\cos (t)+i \sin (t)): 0 \leq t \leq 2 \pi) \geq 0
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$$

In other word, we are minimizing

$$
\begin{aligned}
F(t)= & \left(9 u_{1}-1 u_{2}-18 u_{0}\right) \cos (t) \\
& +\left(8 u_{1}-2 u_{2}-7 u_{0}\right) \sin (t) \\
& -9 u_{1}+2 u_{2}+18 u_{0} \\
= & A \cos (t)+B \sin (t)+C
\end{aligned}
$$

## Example 2

We just give some plots:


Here $(x, y)$-axes are $\left(u_{1}, u_{2}\right)$.

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